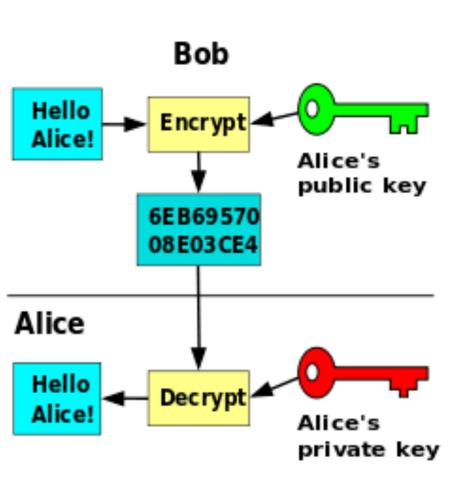
#### In Class Exercise

- Goal: Design a system in which
  - Individuals have sensitive personal data set of attributes (medical records)
  - Data is somehow encrypted by the individual and stored at the cloud
  - A third-party wants to do computation on the data (medical center)
  - The third party also has secret inputs and does not want to share those with the cloud
  - Ideally, user is not involved

### Paillier Cryosystem



- The public key:  $(n, g, h = g^x)$
- Secret key:  $x \in [1, n^2/2]$
- Strong secret:

Factorization of n = zy (z, y are safe primes)

# Paillier Cryptosystem Encryption

- To encrypt a message m ∈ Z\_n
  - -Select a random r ∈ [1, n/4]
  - Generate the ciphertext pair (C1,C2) such that
  - $-C1 = g^r \mod n^2$
  - $-C2 = h^r(1 + mn) \mod n^2$

-[m]=(C1,C2)

# Paillier Cryptosystem Decryption

- The message m can be recovered from [m]=(C1,C2) as follows:
  - $-m = Delta(C2 /C1^x)$
  - $-Delta(u) = [(u-1) \mod n^2]/n$ 
    - For all  $u \in \{u < n^2 \mid u = 1 \mod n\}$

# Paillier Cryptosystem Threshold Encryption

 Assume we randomly split the secret key in two shares x1 and x2,

$$- x = x1 + x2$$

 The Paillier cryptosystem enables an encrypted message (C1,C2) to be partially decrypted to a ciphertext pair (C~1,C~2) using x1 as

```
- C^{1} = C1
```

$$- C^2 = C2 /C1^(x1) \mod n^2$$

• Then, (C~1,C~2) can be decrypted using x2

## Homomorphism

 The product of two ciphertexts is equal to the encryption of the sum of their corresponding plaintexts

 A ciphertext raised to a constant number is equal to the encryption of the product of the corresponding plaintext and the constant

#### **Tasks**

- Decide on the system model and parties involved
- Decide on the threat model for all parties involved
- Design the system
  - Initialization: Key generation, key management, encryption
  - Application: SMC
- Comment on the functions that can be supported
- Comment on the security/privacy of the system
- Comment on the performance
- Comment on the user-friendliness

# System Model







#### **Threat Model**

Semi-honest adversary vs.
 Malicious adversary

 Polynomial-time adversary vs. computationally unbounded adversary

Collusion





## Requirements

- Types of supported queries:
  - Weighted Average
  - Multiplication of ciphertexts
  - Division
  - Comparison/Classification
- Access Control
- Access Patterns

## Design

Initialization

Application(s)